



# Reliably Determining the Outcome of Computer Network Attacks

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## Introduction



- Research Motivation
- Determining Attack Outcome
- IDS Analyst Evasion
- Forging Responses
- Determining Trust
- Conclusion





- Network Intrusion Detection Systems (NIDSs) are more like "attack" detection systems
- Buffer overflow attacks are widespread
- Manual checking of alerts is time consuming and error prone
- Network analysts either overly trust network data or are too paranoid





NIDS detects that an attack is in progress

Decides if the attack is a success or failure



## **Success or Failure?**



- Immediate
  - The intruder makes it obvious
  - Server response to attack
  - Network understanding/mapping
  - Active verification

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			Þ
a2 7a 00 0c 10 00 80 06	29 e1 dd fc 08 00 e3 e7 0a 01 01 30	45 00	
)a 7f db 78 )0 00 4d 69 )f 77 73 20 }5 2e 31 2e	17 c9 c7 09 43 b1 63 72 6f 73 6f 66	50 18	\xC.P.



## **Success or Failure?**



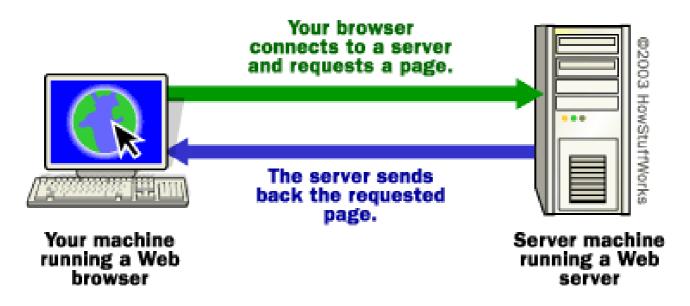
- Delayed
  - Check patches or logs
  - Backdoor signatures
  - Anomaly Detection Traffic analysis/Data Mining



## **Network Traffic Analysis**



# Graphical depiction of a typical request and response



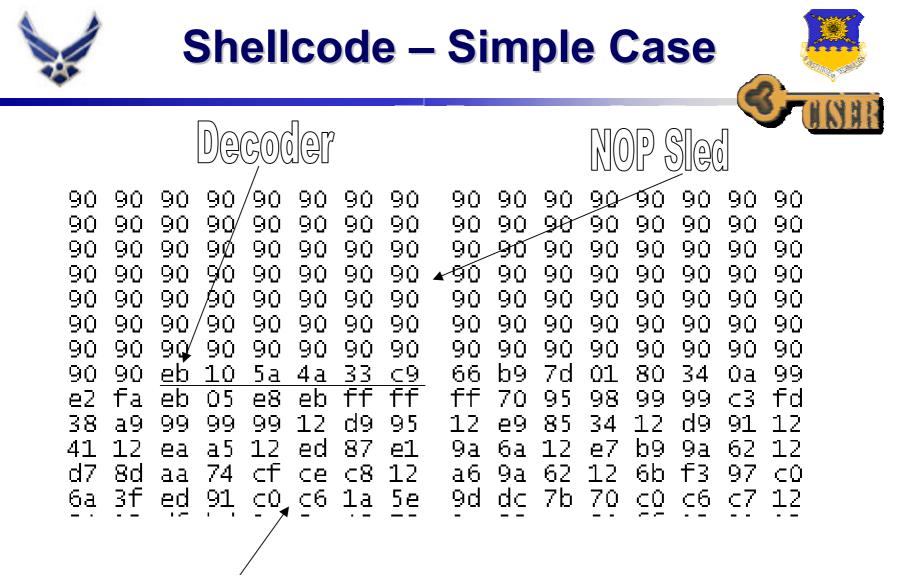


## **Network Traffic Analysis**



## What the NIDS analyst sees

No 1	Source	Destination	Protocol	Info
3 (	10.1.1.10	10.1.1.55	TCP	1346 > http [SYN] Seq=0 Ac
4	10.1.1.55	10.1.1.10	ТСР	http > 1346 [SYN, ACK] Seq
5	10.1.1.10	10.1.1.55	ТСР	1346 > http [ACK] Seq=1 Ac
6	10.1.1.10	10.1.1.55	HTTP	GET /_vti_bin/_vti_aut/fp3
7	10.1.1.55	10.1.1.10	HTTP	HTTP/1.1 500 Server Error
8	10.1.1.55	10.1.1.10	ТСР	http > 1346 [FIN, ACK] Seq
9	10.1.1.10	10.1.1.55	ТСР	1346 > http [ACK] Seq=68 A
10	10.1.1.10	10.1.1.55	ТСР	1346 > http [FIN, ACK] Seq
11	10.1.1.55	10.1.1.10	ТСР	http > 1346 [ACK] Seq=260



She



## **Real World Advice**



- Vendor IDS Signature Guidance
  - "Also look for the result returned by the server. An error message probably indicates the attack failed. If successful, you may see not more traffic in this session (indicating a shell on another port) or non ftpspecific commands being issued"
- Intrusion Signatures and Analysis, Book
  - "The DNS software should be reviewed to ensure that the system is running the latest version"



## **Real World Advice**



- Snort User's Group
  - "In a large number of cases there is nothing preventing the attacker from having the service return the same response as a non vulnerable service"
- IDS User's Group
  - "You still need a trained analyst who knows what the data means to be able to determine what has to be done with it"



## **Real World Advice**



### IDS User's Group

- "In general it's impossible to determine the success of attacks with only a network IDS (NIDS)"
- "For attack like Nimda, you need to check the HTTP response code and see if it return the interesting stuff. For DoS attack, you need to check if the server is crash which will not send back the response"
- "The behavior to that of a non-vulnerable system to an attack is often different and well-defined ..... and there are evasive measures attackers could use to avoid the appearance of success"



# **Test Methodology**



- Experimental Design
  - Windows XP attack system running Ethereal
  - Metasploit Framework used to test/develop exploits
    - Eight buffer overflow vulnerabilities fully tested
  - Windows XP VMWare host running Windows 2000 Server SP 0-4 and Windows XP SP 0-1

## NIDS Test Design

- Vary shellcode Exit Function, test patched and unpatched servers
- Direct measurement of server response, five second captures
- At least three repetitions
- Ensure the vulnerability is tested and not the exploit
- Use VMWare's "Revert to Snapshot" feature



## **Server Response Results**



Exploit	MS Bulletin	Patched Server Response	Un patche d Reponse	Size (bytes)
Apache Chunked	N/A	HTTP/1.1 400 Bad Request	None	542
IIS_WebDAV	03-07	HTTP/1.1 400 Bad Request	None	235
IIS_Nsiislog	03-19/03-22	HTTP/1.1 400 Bad Request	None or 500 Server Error	111
IIS_Printer	01-23	None	None	N/A
IIS_Fp30Reg	03-51	HTTP/1.1 500 Server Error	None	258/261
LSASS	04-11	WinXP: DCERPC Fault Win2K: LSA-DS Response	None	WinXP:92 Win2K:108
RPC DCOM	03-26	RemoteActivation Response	None	92

#### •Is it really this easy?

•Exploit vector, bad input, custom error pages





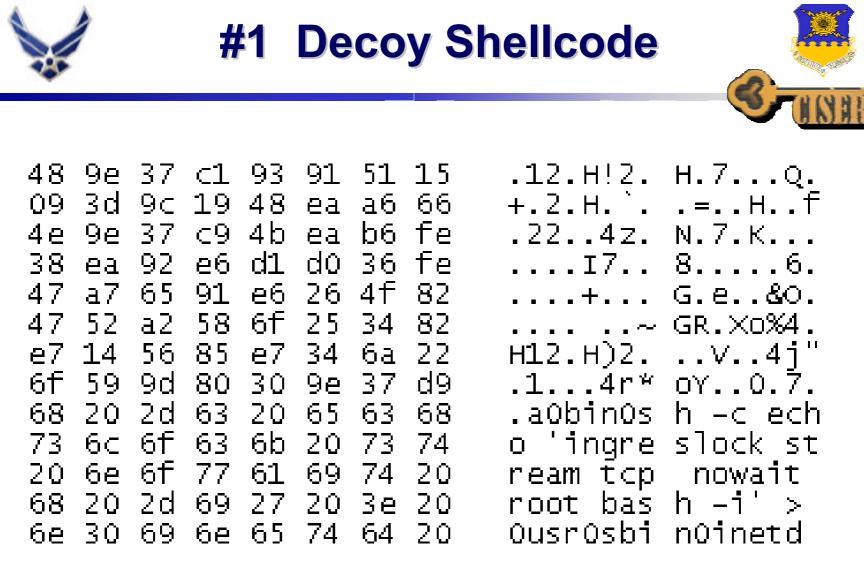


- Typically refers to techniques that evade or disrupt the computer component of the NIDS
- Insertion, Evasion, Denial of Service (DOS)
- Polymorphic shellcode
  - ADMmutate, substitute NOPs
- Mimicry attacks
  - Modify exploit to mimic something else
- NIDS analyst evasion
  - Convince analyst that successful attack has failed





- Training: Analysts recognize UNIX vs. Windows shellcode
- Attack: Create decoy shellcode that appears to target UNIX (i.e. /bin/sh or /etc/inetd.conf), but instead creates a Windows backdoor
- Result: Analyst believes that the attack targets the wrong Operating System







- Training: Analysts look for signs that an intruder could not connect to the backdoor
- Attack: Create shellcode that adds a new user and then send SYN packets to a fake backdoor (i.e., 1524 ingreslock)
- Result: The response from the victim server (RST/ACK) seems to indicate the attack failed



## **#2 Fake Backdoor**



No	Source	Destination	Protocol	Info
1	10.1.1.10	10.1.1.60	тср	1268 > 4444 [SYN] Seq=0
2	10.1.1.60	10.1.1.10	TCP	4444 > 1268 [RST, ACK]
3	10.1.1.10	10.1.1.60	тср	1268 > 4444 [SYN] Seq=0
	10.1.1.60	10.1.1.10	ТСР	4444 > 1268 [RST, ACK]
5	10.1.1.10	10.1.1.60	TCP	1268 > 4444 [SYN] Seq=0
6	10.1.1.60	10.1.1.10	тср	4444 > 1268 [RST, ACK]





- Training: Analysts trust success and failure error codes/characteristics
- Attack: Forge the server response to return the error the analyst is expecting (i.e., HTTP/1.1 400 Bad Request)
- Result: The attack is believed to have failed since the server clearly processed and denied the attack



## **#3 Forged Response**



Jo.+	Source	Destination	Protocol	Info
1	10.1.1.10	10.1.1.55	TCP	1158 > http [SYN] Seq=0 Acl
2	10.1.1.55	10.1.1.10	TCP	http > 1158 [SYN, ACK] Seq
3	10.1.1.10	10.1.1.55	ТСР	1158 > http [ACK] Seq=1 Ack
4	10.1.1.10	10.1.1.55	HTTP	Continuation
- 5	10.1.1.55	10.1.1.10	HTTP	HTTP/1.1 400 Bad Request
6	10.1.1.10	10.1.1.55	ТСР	1158 > http [АСК] Seq=1304
7	10.1.1.55	10.1.1.10	TCP	http > 1158 [RST] Seq=91 A



- Find the socket descriptor associated with the attacker's connection
- Findsock
  - Use getpeername and attacker's source port
  - Doesn't work through NAT/proxies
- Findtag
  - Use *ioctlsocket* and *FIONREAD* to read in a hardcoded tag
  - Requires an additional packet after overflow



## **Findtag and Findsock**



Source	Destination	Protocol	Info
10.1.1.10	10.1.1.55	TCP	1586 > http [SYN] Seq=0 /
10.1.1.55	10.1.1.10	ТСР	http > 1586 [SYN, ACK] Se
10.1.1.10	10.1.1.55	TCP	1586 > http [ACK] Seq=1 /
10.1.1.10	10.1.1.55	HTTP	Continuation
10.1.1.55	10.1.1.10	TCP	http > 1586 [ACK] Seq=1 /
10.1.1.10	10.1.1.55	HTTP	Continuation
10.1.1.55	10.1.1.10	ТСР	http > 1586 [ACK] Seq=1 /
10.1.1.55	10.1.1.10	HTTP	HTTP/1.1 400 Bad Request
10.1.1.10	10.1.1.55	ТСР	1586 > http [ACK] Seq=13(
10.1.1.55	10.1.1.10	ТСР	http > 1586 [RST] Seq=91
09 80 00 11		8 00 45	
40 00 80 06	78 9e 0a 01 0		
00 50 2b 60	54 ce 57 eb 5	b 8C 50	
00 00 6d 73	66 21		f!

Hard-coded: 40 bytes Universal: 90 bytes

Process Injection (minimum API calls): 255 bytes







- Create the packet from scratch using raw sockets (Windows 2000, XP, 2003 targets)
- Rawsock
  - Socket, setsockopt, sendto
  - Requires administrative privilege
  - Requires that attacker capture Initial Sequence Numbers and calculate checksum
  - Hardcoded: 350 bytes





- Use techniques introduced in public exploits to locate the connection ID during overflows in Internet Server API (ISAPI) extensions
  - Locate Extension Control Block
  - Find connection ID (socket handle equivalent)
  - Pick default error message (ServerSupportFunction Send Response Header)
  - Send forged message (Writeclient)
- Smaller shellcode, does not rely on the error message size (unless custom page)



# **Server Response Trust**



- Payload Size Analysis
  - Calculate payload size and compare to minimum forging requirements. In most cases at least 350 bytes is required for forging and backdoor
- Check if shellcode is known
  - Match shellcode to common exploits available on the internet (an automated tool would be best)
  - Keep database of most used exploits/payloads
- Decode the shellcode to determine function
  - Requires expert skill or sophisticated computer program







Source	Destination	Protocol	Info
10.1.1.10	10.1.1.55	TCP	2123 > http [SYN] Seq=0 Ack=0 Win=64
10.1.1.55	10.1.1.10	TCP	http > 2123 [SYN, ACK] Seq=0 Ack=1 W
10.1.1.10	10.1.1.55	ТСР	2123 > http [АСК] Seq=1 Ack=1 Win=64
10.1.1.10	10.1.1.55	HTTP	POST /scripts/nsiislog.dll HTTP/1.1
10.1.1.55	10.1.1.10	TCP	http > 2123 [ACK] Seq=1 Ack=2921
10.1.1.10	10.1.1.55	HTTP	Continuation
10.1.1.55	10.1.1.10	HTTP	HTTP/1.1 100 Continue
10.1.1.55	10.1.1.10	TCP	http > 2123 [АСК] Seq=90 Ack=7301 Wi
10.1.1.10	10.1.1.55	HTTP	Continuation
10.1.1.55	10.1.1.10	ТСР	http > 2123 [ACK] Seq=90 Ack=58401 win=:
10.1.1.55	10.1.1.10	TCP	http > 2123 [ACK] Seq=90 Ack=61321 Win=:
10.1.1.55	10.1.1.10	TCP	http > 2123 [ACK] Seq=90 Ack=64241 Win=:
10.1.1.55	10.1.1.10	TCP	http > 2123 [ACK] Seq=90 Ack=65916 Win=
10.1.1.55	10.1.1.10	TCP	[TCP Dup ACK 44#1] http > 2123 [ACK] Se
10.1.1.55	10.1.1.10	HTTP	HTTP/1.1 400 Bad Request
10.1.1.55	10.1.1.10	ТСР	http > 2123 [FIN, ACK] Seq=201 Ack=6591
10.1.1.10	10.1.1.55	ТСР	2123 > http [ACK] Seq=65916 Ack=202 Win

## Success or failure?







07b0 90 90 90 90 90 90 90 90 90 90 90 90 90 90 90 90 07c0 90 90 90 90 d9 ee d9 74 24 f4 5b 31 c9 b1 2d 81 .....t **\$**.[1..-. 07d0 73 17 d3 18 83 eb fc e2 06 f2 d3 ee a4 f4 18 ee 2f bf f0 07e0 f7 86 4e b9 3c f6 2f 96 24 65 d6 60 ef 07f0 2f 30 2f 4e 58 52 f6 89 38 ef 4f 2a a7 e7 93 ef e2 95 1f0800 4 a e7 12 13 b4 d6 44 a7 2f 6b de 19 29 0810 4f 21 80 с7 6d Of 2f 89 4f b5 93 23 92 3c ef e2 0820 ef 58 f2 38 93 2f e2 d8 42 a 5 ea d2 f0 05 a6 b1 8e 0830 fa 43 2c ac d3 46 84 94 8a 65 bd 58 43 e2 7c bf 3f 0840 2f 88 65 58 43 f7 bb 96 38 -04 e6 a0 aa e7 0850 f3 f6 f7 bb 80 10 dd 5b 4b 2d 14 2e e7 38 03 86 72 05 03 3e d3 f0 0860 4b 1df8 8e 5b 5b e7 0f ee се 0870 2c e7 8d b7 36 8b C1 f3 68 5b c9 dc b6 38 c7 87 fd 0880 b4 38 df 94 29 <0 95 fd 2d de d3 18 ca a4 4d 0890 4d 4d 4d 4d 4d 4 d 4d 08a0 4d 4 d4d 4d 4d 4d 4d

s.... ..N./.<. /.\$e.. NXR./.8. 00\*./... D.../k...) ].... 0!#...m. /.<.0... .×B..8.. 1. . ..C,..F. ..|e.×⊂. /..e.×⊂. .8[.ĸ. к...Г.г. ..6. 8. . -..h. MMMMMMM MMMMMMMM MMMMMMMM MMMMMMMMM

Payload size = 088e - 07c4 = CA (hex) = 202 bytes

## Is forging possible?







Source	Destination	Protocol	Info
10.1.1.10	10.1.1.55	TCP	1170 > epmap [SYN] Seq=0 Ack=(
10.1.1.55	10.1.1.10	ТСР	epmap > 1170 [SYN, ACK] Seq=0
10.1.1.10	10.1.1.55	TCP	1170 > epmap [ACK] Seq=1 Ack=1
10.1.1.10	10.1.1.55	DCERPC	Bind: call_id: O UUID: REMACT
10.1.1.55	10.1.1.10	DCERPC	Bind_ack: call_id: 0 accept ma
10.1.1.10	10.1.1.55	REMACT	RemoteActivation request
10.1.1.55	10.1.1.10	TCP	epmap > 1170 [ACK] Seq=61 Ack:
10.1.1.10	10.1.1.55	TCP	1170 > epmap [ACK] Seq=1747 A
10.1.1.10	10.1.1.55	TCP	1177 > 4444 [SYN] Seq=0 Ack=0
10.1.1.55	10.1.1.10	TCP	4444 > 1177 [RST, ACK] Seq=0 /
10.1.1.10	10.1.1.55	TCP	1177 > 4444 [SYN] Seq=0 Ack=0
10.1.1.55	10.1.1.10	TCP	4444 > 1177 [RST, ACK] Seq=0 /
10.1.1.10	10.1.1.55	TCP	1177 > 4444 [SYN] Seq=0 Ack=0
10.1.1.55	10.1.1.10	TCP	4444 > 1177 [RST, ACK] Seq=0 /
10.1.1.10	10.1.1.55	ТСР	1170 > epmap [FIN, ACK] Seq=1.

## Success or failure?







## Success or failure?







**Payload Database** 

**Attacker's Shellcode** 

Do they match?



## What about Linux?



- Server Response Characteristics
- Forging attacks
- Trust Determination







- The outcome of many buffer overflow attacks can be automatically determined based on network data alone
- There is no difference between a forged and a legitimate response
  - However it can be determined, in most cases, if forging is possible
- NIDS developers should leave as little to the analyst as possible (obvious, but more needs to be done)
  - When possible block malicious traffic
  - Post-processing of response/validity calculation







### **Questions?**

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